

Context Matching for Ambient Intelligence Applications

Andrei Olaru

University Politehnica of Bucharest

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- Conclusion

Context Matching for Ambient Intelligence Applications

overview



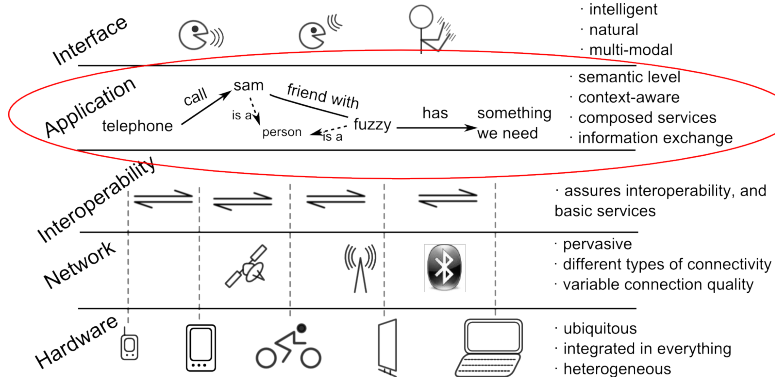
Ambient Intelligence – or Aml – is a ubiquitous electronic environment that supports people in their daily tasks, in a proactive, but "invisible" and non-intrusive manner. [Ducatel et al., 2001]

- ▶ We can view an Aml environment as a system of "information conveyers" [Weiser, 1993]
- ▶ **Software agents** are an appropriate implementation for Aml systems [Ramos et al., 2008]

· The ideal features of Aml are also its greatest challenges:

- ▶ Uniformity / unification
- ▶ Scalability
- ▶ Availability / reliability

Our approach: Build a multi-agent system for the context-aware exchange of information in an Aml environment – the [AmIciTy initiative](#). [Olaru et al., 2013]



Aml Layers (based on [El Fallah Seghrouchni, 2008])

Context is any information that can be used to characterize the situation of entities (i.e. a person, place or object) that are considered relevant to the interaction between a user and an application, including the user and the application themselves. [Dey, 2001]

· Example of context-aware scenario:

If I am passing near my bank, during working hours, but I am not currently walking together with someone, I want to be reminded to go to the bank.



Context is any information that can be used to characterize the situation of entities (i.e. a person, place or object) that are considered relevant to the interaction between a user and an application, including the user and the application themselves. [Dey, 2001]

· Example of context-aware scenario:

*If I am passing **near my bank**, **during working hours**, but I am not currently walking together with someone, I want to be reminded to go to the bank.*

location *time*

social



Context is any information that can be used to characterize the situation of entities (i.e. a person, place or object) that are considered relevant to the interaction between a user and an application, including the user and the application themselves. [Dey, 2001]

- Example of context-aware scenario:

or



Context is any information that can be used to characterize the situation of entities (i.e. a person, place or object) that are considered relevant to the interaction between a user and an application, including the user and the application themselves. [Dey, 2001]

· Example of context-aware scenario:

Having received an email, I want the Aml system to detect if it is a call for papers and to notify me if I haven't sent a paper .



Context is any information that can be used to characterize the situation of entities (i.e. a person, place or object) that are considered relevant to the interaction between a user and an application, including the user and the application themselves. [Dey, 2001]

- Example of context-aware scenario:

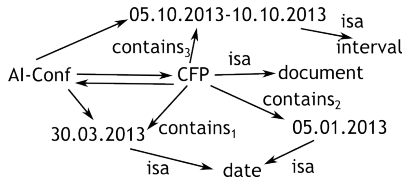
*Having received an email, I want the Aml system to detect if **it is a call for papers** association and to notify me **if I haven't sent a paper** association .*



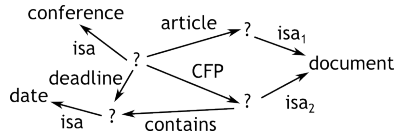
We define **context matching** as matching *context patterns* against the current *context graph*.

[Olaru et al., 2011]

- the context graph represents relations between concepts;
- context patterns are graphs featuring generic nodes;



context graph



context pattern



- Context matching is used for
 - ▶ knowledge integration → incoming information matches the agent's patterns; ← perceiving
 - ▶ situation recognition → the pattern matches part of the context graph; ← reactivity
 - ▶ problem detection → only part of the pattern matches the context; ← pro-activity / anticipation
 - ▶ sharing information → other agent's patterns match the context graph. ← cooperation

- Local matching helps scalability and privacy-awareness. ← reasoning and detection are performed locally.



Problem statement: devise an algorithm that makes context matching (underpinned by graph matching) a valid approach for the implementation of a context-awareness mechanism in agents that reside on devices a various sizes.

- that is, an algorithm that is tractable for cases specific to our problem:
 - ▶ graphs have mostly labeled edges;
 - ▶ there may be a reasonable amount of generic nodes in graph patterns;
 - ▶ the size of the context graph and context patterns will be adequate to the capabilities of the device;

Existing graph matching algorithms date from the 70's to present times

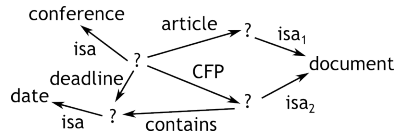
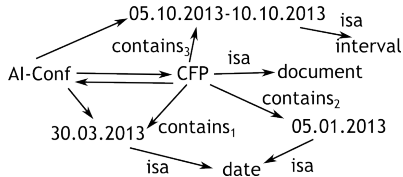
[Cordella et al., 2004]

- ▶ exact vs. inexact matching;
- ▶ traditional algorithms match unlabeled, undirected graphs → **modifications are needed**;
- ▶ studied algorithms:
 - McGregor – exploring the entire state space; [McGregor, 1982]
 - Bron-Kerbosch, Durand-Pasari, Akkoyunlu and Balas-Yu – searching maximal cliques in the associations graph;

[Bron and Kerbosch, 1973, Akkoyunlu, 1973, Balas and Yu, 1986, Durand et al., 1999]
 - Koch – searching maximal cliques in the modular product of the edges;

[Koch, 2001]
 - Larrosa – modeling the matching problem as CSP. [Larrosa and Valiente, 2002]
- ▶ an adaptation of various algorithms has been implemented and comparison has been performed. [Dobrescu and Olaru, 2013]





Example

Model

Context graph

$$CG_A = (V, E)$$

$$V \subseteq \text{Concepts}$$

$$E = \{ \text{edge}(\text{from}, \text{to}, \text{value}) \mid \\ \text{from}, \text{to} \in V, \\ \text{value} \in \text{Relations} \}$$

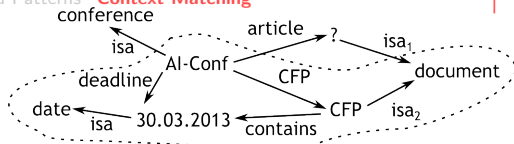
Context pattern

$$G_s^P = (V_s^P, E_s^P)$$

$$V_s^P \subseteq \text{Concepts} \cup \{?\}$$

$$E_s^P = \{ \text{edge}(\text{from}, \text{to}, \text{value}) \mid \\ \text{from}, \text{to} \in V_s^P, \\ \text{value} \in \text{Relations} \cup \{\lambda\} \}$$





Example

Model

Match: $M_{A-si}(G'_A, G_m^P, G_x^P, f_v, k)$

$G'_A \subseteq CG_A$, $G_m^P = (V_m^P, E_m^P) \subseteq G_s^P$ – matched subgraph, pattern solved part

$G_x^P = (V_x^P, E_x^P) \subseteq G_s^P$ – pattern unsolved part

$G_m^P \cup G_x^P = G_s^P$, $V_m^P \cap V_x^P = E_m^P \cap E_x^P = \emptyset$ – no solved & unsolved intersection

$f_v : V_s^P \rightarrow V'$ – vertex correspondence (bijective) – with:

- $\forall v^P \in V_m^P, v^P = ?$ or $v^P = f(v^P)$
- $\forall \text{edge}(v_i^P, v_j^P, \text{value}) \in E_m^P, \text{edge}(f(v_i^P), f(v_j^P), \text{value}) \in E'$



- ▶ Start from all valid matches of one edge in the pattern with one edge in the graph;
- ▶ For each initial match, detect which other matches are valid merger candidates;
- ▶ Iterate over matches and create new matches, by merging them to their merger candidates;



- a match is represented as $M(G', G_m^P, f_v, f_e, fr, MC, MO, k)$

$G' \subseteq CG_A$ – matched subgraph of CG_A

$G_m^P \subseteq G_s^P$ – matched part of the pattern

$f_v : V_m^P \rightarrow V'$ – node function (bijective)

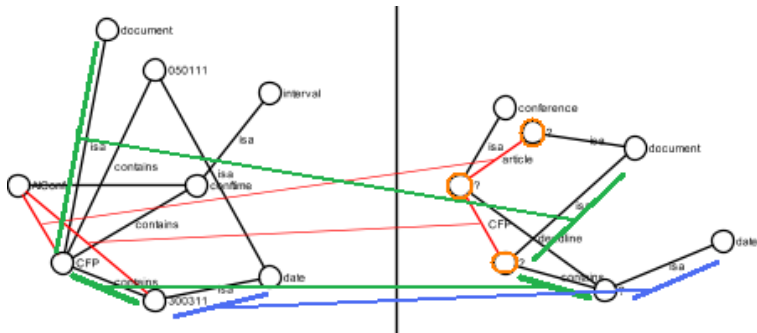
$f_e : E_m^P \rightarrow E'$ – edge function (bijective)

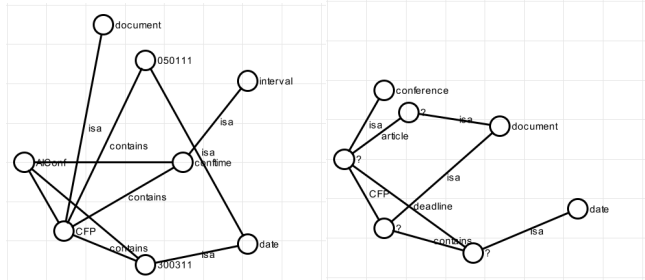
$fr \subseteq V_m^P$ – frontier

MC – merger candidates

MO – 'outer' merger candidates

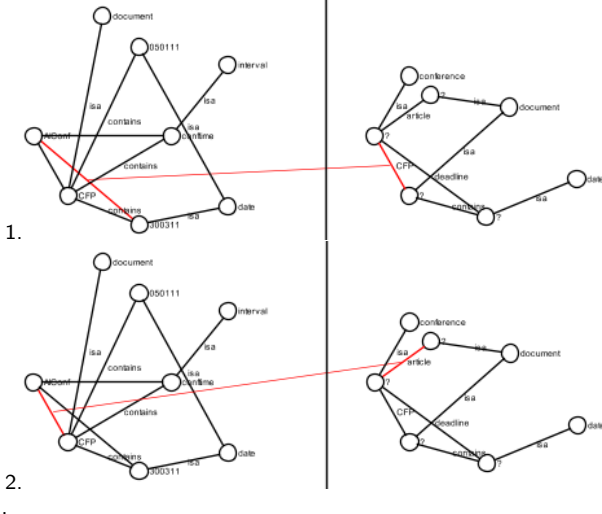
$k = |E_s^P| - |E_m^P|$ – missing edges



$Match(G, G^P)$


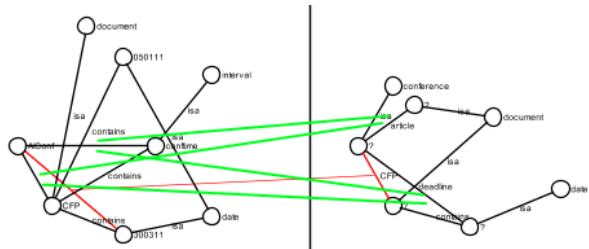
$Match(G, G^P)$

AddInitialMatches



$Match(G, G^P)$
AddInitialMatches

Immediate candidates

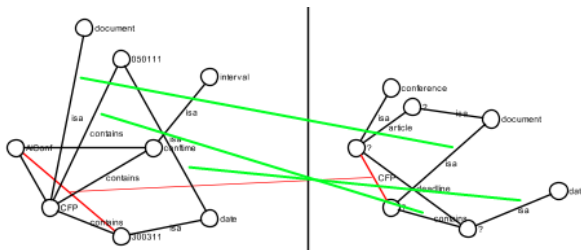


$Match(G, G^P)$

AddInitialMatches

Immediate candidates

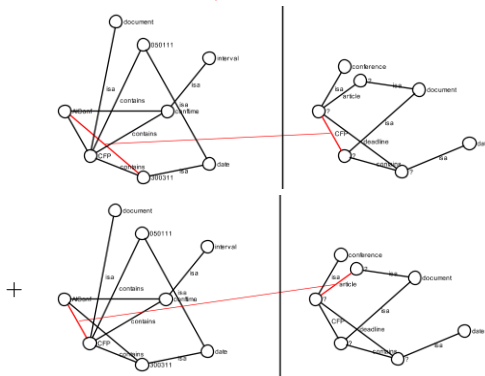
Outer candidates



$Match(G, G^P)$
AddInitialMatches

Immediate candidates

Outer candidates

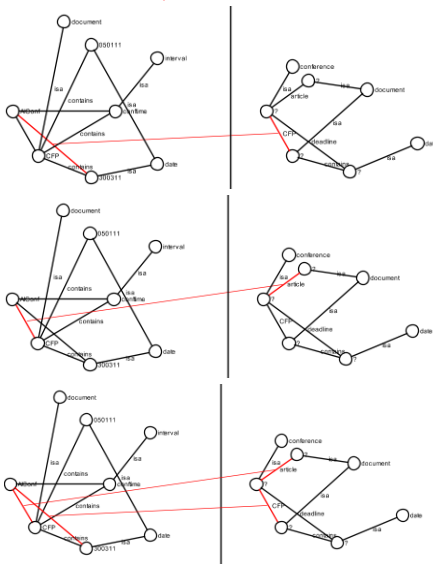
 for each M', M''


$Match(G, G^P)$
 $AddInitialMatches$

Immediate candidates

Outer candidates

for each M', M''

 $Merge(M', M'')$


for each $(e_{kp}^P, e_{kg}) \in E^P \times E$

if e_{kp}^P and e_{kg} match

AddInitialMatches

create new initial single-edge match M

with $E_m^P = \{e_{kp}^P\}$ and $E' = \{e_{kg}\}$

search *MatchQueue*

for all match candidates for M

add M to *MatchQueue*

Match(G, G^P)

for each $M' \in \text{MatchQueue}$, for each $M'' \in M'.MC$

remove M'' from $M'.MC$ and M' from $M''.MC$

GrowMatches

Merge

(M', M'')

$$V = V' \cup V''; E = E' \cup E''$$

$$V_m^P = V_m^{P'} \cup V_m^{P''}; E_m^P = E_m^{P'} \cup E_m^{P''}$$

$$f_v = f_v' \cup f_v''; f_e = f_e' \cup f_e''$$

$$fr = \{v^P \in fr' \cup fr'' \mid \exists e^P \text{ adj } v^P, e^P \notin E_m^P\}$$

$$MC = (MC' \cap MC'') \cup (MC' \cap MO'') \\ \cup (MC'' \cap MO')$$

$$MO = MO' \cap MO''$$



- While the classic problem of matching undirected, unlabeled graphs is NP-complete, for the problem at hand the algorithm behaves significantly better.

AddInitialMatches Creates a maximum of $m \times m^P$ matches, with many less if edges are labeled.

In our example ($m = 11$, $m^P = 8$) there are 19 initial matches.

Each initial match is tested against the other matches for compatibility.

Complexity: $\mathcal{O}(m \times m^P) + \mathcal{O}(\text{initialMatches}^2)$

Merge(M' , M'')

Adds all edges and nodes to the new match

Merges immediate and outer merger candidates

Complexity: $\mathcal{O}(|E_m^{P'}| + |E_m^{P''}|)$

GrowMatches

Iterates over the match queue and merges matches with their candidates.

Complexity:

$\mathcal{O}(\text{initialMatches} \cdot \log(\text{initialMatches}) \cdot \text{average } |E_m^P|)$



Algorithm Evaluation

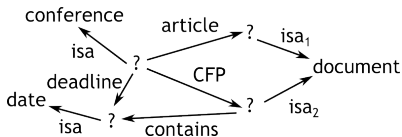
Algorithm	<i>Akkoyunlu</i>	<i>Bron-Kerbosch</i>	<i>Balas-Yu</i>	<i>Durand-Pasari</i>	<i>Our algorithm</i>
Expanded edges:					
Small example	124	120	135	119	34
Initial example	5431	5440	6423	5219	2459
No labeled edges	7054	9454	15843	9060	7581
No labels	326044	371943	578401	367725	108902
Large example	20470	19989	22170	18322	11834

Comments on algorithms not in the table:

- ▶ McGregor expands less edges, but many more nodes;
- ▶ Larossa expands significantly less edges, but can only provide full pattern matches.



- **Text-based representation** for directed graphs that is easy to read by humans.
 - it relies on building a tree of paths, starting with the longest path.



manual

?#1
 $(\xrightarrow{CFP} ?\#4$
 $(\xrightarrow{contains} ?\#3 \xrightarrow{isa} date)$
 $\xrightarrow{isa} document)$
 $(\xrightarrow{deadline} *?\#3)$
 $(\xrightarrow{article} ?\#2 \xrightarrow{isa} *document)$
 $\xrightarrow{isa} conference$

textual

```

merging match [-] (k=6): AIconf (->CFP) ->300311 : ?#3 (-article->?#2) -CFP->?#4
and match [3:2] (k=7): AIconf->conftime : ?#3-deadline->?#2
new match: match [-] (k=5): AIconf (->CFP) (->300311) ->conftime : ?#5
(-article->?#4) (-CFP->?#6) -deadline->?#2
  
```

Console (ASCII)



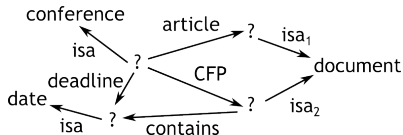
Textual

Graphical

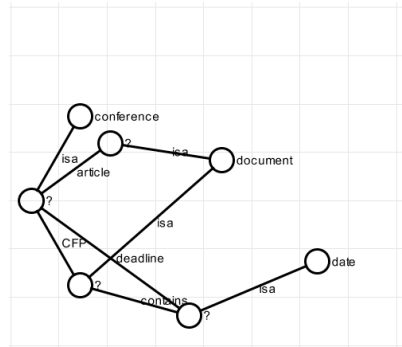
Visualization Tools

► Graphical representation for directed graphs

- relies on the textual representation to build paths with a low number of links between paths
- lays nodes out on concentric 120° arcs



manual



automatic layout



Conclusion and Future Work

- ▶ We have developed an efficient algorithm for the partial matching of context patterns against context graphs.
 - ▶ It relies on creating all valid single-edge matches and then growing matches by merging.
 - ▶ The algorithm has been implemented and it has been compared with other, traditional, graph matching algorithms.
- Future work:
- ▶ Further comparison with other algorithms using automatic graph generation and testing tools.
 - ▶ Integration context matching as reasoning and decision engine in an agent-based platform for Ambient Intelligence.



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Thank You!

Any Questions?





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